In the Specification

Please replace the last paragraph of page 15 (starting at line 29 and continuing on page 16, line 17) with the following:

Consider frames received by a switch 150 of the present invention from computer 120 and intended for RAID system 106 N_Port 118. The headers of each of these frames are decoded by switch 150. In the network as illustrated, frames having D_ID field 3 corresponding to a destination of N_Port 118 may reach that destination through a path through switches 152 and 154, and through a second path through switch 154 directly. A hash function of the DID field 300 and at least one bit of the OX_ID field 304 of the header are therefore used to index routing table 180 to select the outgoing switch port. The routing table 180 has the structure illustrated in Figures 3C and 3D 4C or 4D. The hash function is selected such that all entries of the routing table 180 that may be selected by a valid D_ID field 300 correspond to a valid outgoing port on a path to the N_Port identified by D_ID that is distinct from the incoming switch port.

Please replace the last paragraph of page 19 (starting at line 18 and continuing on page 20, line 2) with the following:

Routing tables of the present invention may be implemented in firmware or hardware of the switch. It is known that implementation of routing tables in hardware provides advantage for switches having heavy load and large numbers of switch ports. In a hardware implementation, routing table 350 of Figure 3C, 366 of Figure 3D, 404 of Figure 4A, or 428 of Figure 4B, may be implemented with a static RAM, and the portmap table 432 with a second static

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RAM. In such an embodiment, the routing table address inputs are multiplexed (365) so it can be written by a processor 363 of the switch such that the processor 363 can maintain the routing table. The routing table is thereby addressable by either the address generated by the hash function or functions, or by an address generated by the processor.